

Computer Architecture and Assembly Language

Practical Session 9

Position Independent Code (PIC)
can be executed at
any memory address
without modification

```

section .data
    var: dd 0x10
    str: db "The result is %d",10,0

    extern printf
    global main

section .text
main:
    call func
    push dword [var]
    push str
    call printf
    add esp, 8

    mov eax,1
    int 0x80

func:
    push ebp
    mov ebp, esp
    mov ecx, 3
    calcLoop:
        add dword [var],2
        loop calcLoop

    pop ebp
    ret

```



Executable binary file (partially):

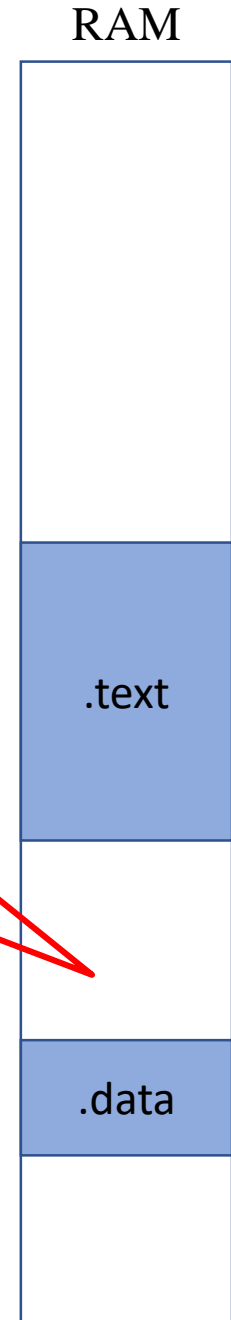
.data	start address	content
.text	start address	content



What if a loader is "confused" and place the sections in wrong position ?

Would the process run correctly ?

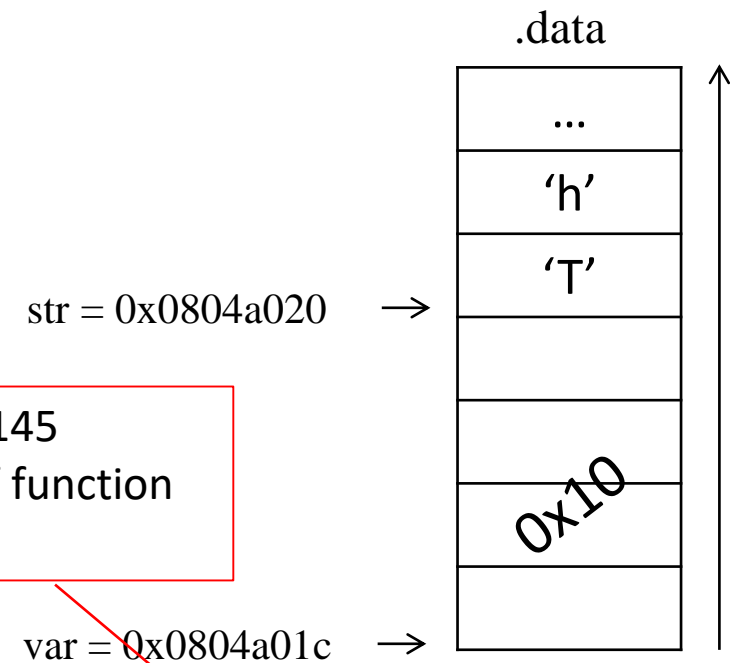
If the program is written in PIC manner, it would run despite of the wrong position. Otherwise, it would not.



```

1      section .data
2 00000000 10000000      var:      dd 0x10
3 00000004 54686520726573756C-      str:      db "The result is %d",10,0
4 0000000D 742069732025640A00
5
6      extern printf
7      global main
8      section .text
9
10     main:
11     call func
12     push dword [var]
13     push str
14     call printf
15     add esp, 8
16
17     mov eax,1
18     int 0x80
19
20     func:
21     push ebp
22     mov ebp, esp
23     mov ecx, 3
24
25     calcLoop:
26     add dword [var],2
27     loop calcLoop
28     pop ebp
29     ret

```



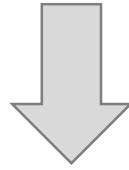
0xffffebb=-0x145
Offset of printf function
in section .text

```

Breakpoint 1, 0x08048410 in main ()
(gdb) x /1xg $eip
0x8048410 <main>: 0x1c35ff0000001ae8
(gdb)
0x8048418 <main+8>: 0x0804a020680804a0
(gdb)
0x8048420 <main+16>: 0x08c483fffffebbe8
(gdb)
0x8048428 <main+24>: 0x5580cd00000001b8
(gdb)
0x8048430 <func+1>: 0x8300000003b9e589
(gdb)
0x8048438 <calcLoop+1>: 0xf7e2020804a01c05
(gdb)
0x8048440 <calcLoop+9>: 0x906690669066c35d

```

No direct usage of labels



- No library functions → only system calls 😞
(their code is not necessarily PIC and thus not safe to use)
- Indirect usage of labels → jump, call, loop commands are PIC 😊
(note that there are several types of jump and call commands)
- Variables address detection in run time → single section code

```
section .data
    var: dd 0x10

section .text

global main
main:
    add dword [var],2

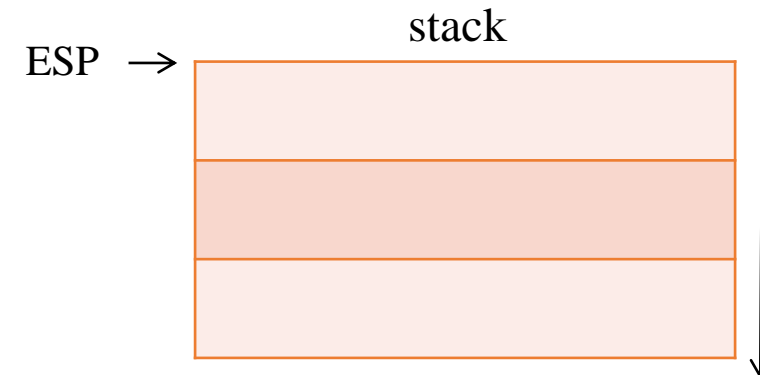
    mov eax,1
    int 0x80
```

```
section .data
    var: dd 0x10

get_my_loc:
    call next_i
next_i:
    pop ecx
    ret

global main
main:
    call get_my_loc
    sub ecx, next_i - var
    add dword [ecx],2

    mov eax,1
    int 0x80
```



```
section .data
    var: dd 0x10

section .text

global main

main:

    add dword [var],2

    mov eax,1
    int 0x80
```

```
section .data
    var: dd 0x10

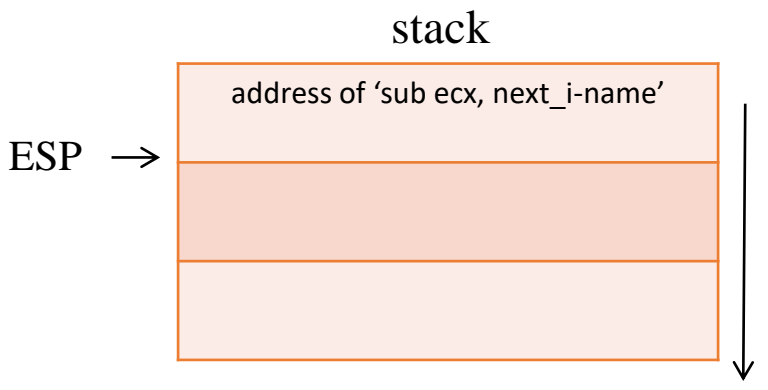
get_my_loc:
    call next_i

next_i:
    pop ecx
    ret

global main

main:
    call get_my_loc
    sub ecx, next_i - var
    add dword [ecx],2

    mov eax,1
    int 0x80
```



```
section .data
    var: dd 0x10

section .text

global main

main:
    add dword [var],2

    mov eax,1
    int 0x80
```

```
section .data
    var: dd 0x10

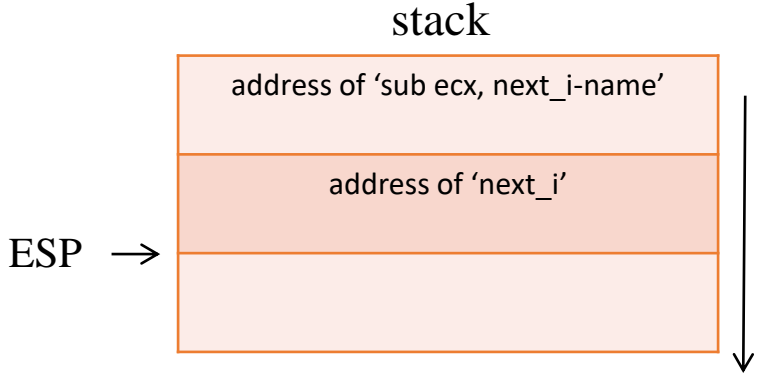
get_my_loc:
    call next_i

next_i:
    pop ecx
    ret

global main

main:
    call get_my_loc
    sub ecx, next_i - var
    add dword [ecx],2

    mov eax,1
    int 0x80
```




```
section .data
    var: dd 0x10

section .text

global main
main:
    add dword [var],2

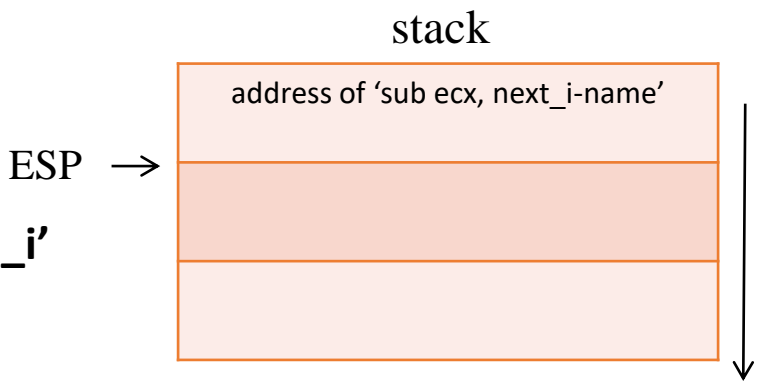
    mov eax,1
    int 0x80
```

```
section .data
    var: dd 0x10

get_my_loc:
    call next_i
next_i:
    pop ecx ; ecx gets address of 'next_i'
    ret

global main
main:
    call get_my_loc
    sub ecx, next_i - var
    add dword [ecx],2

    mov eax,1
    int 0x80
```



```
section .data
    var: dd 0x10

section .text

global main

main:
    add dword [var],2

    mov eax,1
    int 0x80
```

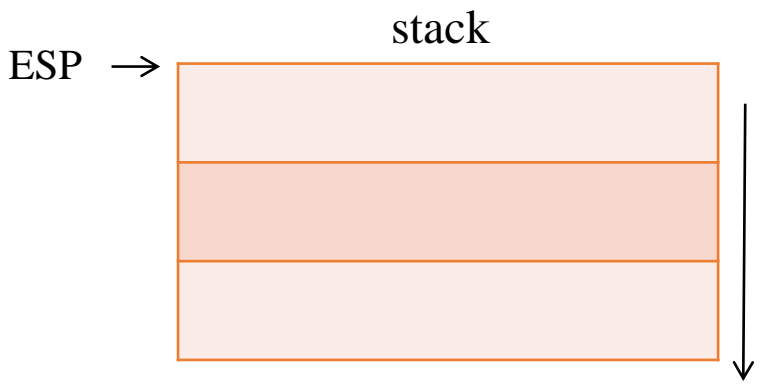
```
section .data
    var: dd 0x10

get_my_loc:
    call next_i
next_i:
    pop ecx
    ret

global main

main:
    call get_my_loc
    sub ecx, next_i - var
    add dword [ecx],2

    mov eax,1
    int 0x80
```



```
section .data
    var: dd 0x10

section .text

global main
main:
    add dword [var],2

    mov eax,1
    int 0x80
```

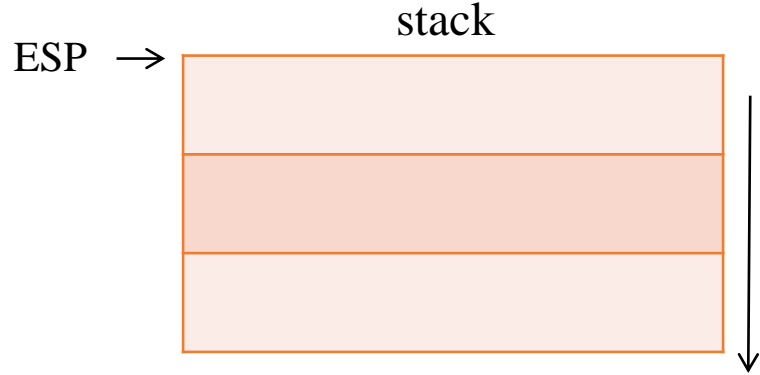
```
section .data
    var: dd 0x10

get_my_loc:
    call next_i
next_i:
    pop ecx
    ret

global main
main:
    call get_my_loc
    sub ecx, next_i - var
    add dword [ecx],2

    mov eax,1
    int 0x80
```

; ecx = 'next_i' address - ('next_i' address - 'var' address)



the address difference between "next_i" and "var" is constant even if the code changes it's position

section .text

```
name:    db      "Hello",10
nameLen  equ     $ - name
```

```
global _start
```

```
get_my_loc:
```

```
call     next_i
```

```
next_i:
```

```
pop      ecx
```

```
ret
```

```
_start:
```

```
call     get_my_loc
```

```
sub      ecx, next_i - name
```

```
mov      edx, nameLen
```

```
mov      eax, 4
```

```
mov      ebx, 1
```

```
int      80h
```

```
mov      eax, 1
```

```
int      80h
```



Does this code
is PIC ?

section .text

```
name:    db      "Hello",10  
nameLen equ     $ - name
```

```
global _start
```

```
get_my_loc:
```

```
call     next_i
```

```
next_i:
```

```
pop     ecx
```

```
ret
```

```
_start:
```

```
call    get_my_loc
```

```
sub     ecx, next_i - name
```

```
mov     edx, nameLen ←
```

```
mov     eax, 4
```

```
mov     ebx, 1
```

```
int     80h
```

```
mov     eax, 1
```

```
int     80h
```

Does this code
is PIC ?

may we use 'nameLen' label directly ?