

## Homework assignment no. 2

1. (a) Draw a polygon  $P$  and place guards in it, such that the guards cover the boundary of  $P$ , but there exists a point in the interior of  $P$  that is not seen by any of the guards.  
(b) A *rectilinear polygon* is a simple polygon whose edges are either horizontal or vertical. Let  $P$  be a rectilinear polygon with  $n$  vertices. Show that  $\lfloor n/4 \rfloor$  guards are sometimes necessary to guard  $P$ .
2. A simple polygon  $P$  is *star-shaped* if there exists a point  $c \in P$ , such that for any point  $p \in P$  the line segment  $\overline{cp}$  is contained in  $P$ . The point  $c$  is called a *center point* of  $P$ .
  - (a) Let  $P$  be a star-shaped polygon, and let  $c$  be a center point of  $P$ . Show that, given a query point  $q$ , one can determine in  $O(\log n)$  time whether  $q$  lies in  $P$ . Assume that  $P$  is given as an array of its  $n$  vertices in sorted order along the boundary.
  - (b) Give an expected linear-time algorithm to decide whether a simple polygon is star-shaped.
3. Given  $n$  inequalities  $a_i x + b_i y \geq 1$ , for  $i = 1, \dots, n$ , describe an expected linear-time algorithm that finds a point  $(x, y)$  (if exists) that (i) satisfies all these inequalities, and (ii) is nearest to the origin under the  $L_1$  distance (where  $d_1(p, q) = |q_x - p_x| + |q_y - p_y|$ ).
4. Let  $R$  be a set of  $n$  axis-parallel rectangles in the plane. We would like to be able to report all rectangles in  $R$  that are *fully* contained in a query axis-parallel rectangle. Describe a data structure of size  $O(n \log^{c_1} n)$  that supports such queries in time  $O(\log^{c_2} n + k)$ , where  $c_1, c_2$  are constants and  $k$  is the number of reported rectangles.
5. **Bonus:** Let  $S$  be a set of  $n$  line segments on the  $x$ -axis. For a segment  $s \in S$ , consider the segments in  $S$  whose left endpoint lies in  $s$  and whose right endpoint lies to the right of  $s$ . Let  $p(s)$  be the rightmost endpoint among the right endpoints of these segments. Describe an algorithm that finds the point  $p(s)$ , for each  $s \in S$ , in total time  $O(n \log n)$ .

**Submission:** December 29, 2015.