

Sparse Graph Spanners

(2005; Elkin, Peleg)

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SYNONYMS: $(1 + \epsilon, \beta)$ -spanners; almost additive spanners.

1 PROBLEM DEFINITION

For a pair of numbers α, β , $\alpha \geq 1$, $\beta \geq 0$, a subgraph $G' = (V, H)$ of an unweighted undirected graph $G = (V, E)$, $H \subseteq E$, is an (α, β) -spanner of G if for every pair of vertices $u, w \in V$, $\text{dist}_{G'}(u, w) \leq \alpha \cdot \text{dist}_G(u, w) + \beta$, where $\text{dist}_G(u, w)$ stands for the distance between u and w in G . It is desirable to show that for every n -vertex graph there exists a sparse (α, β) -spanner with as small values of α and β as possible. The problem is to determine asymptotic tradeoffs between α and β on one hand, and the sparsity of the spanner on the other.

2 KEY RESULTS

The main result of Elkin and Peleg [4] establishes the existence and efficient constructibility of $(1 + \epsilon, \beta)$ -spanners of size $O(\beta n^{1+1/\kappa})$ for every n -vertex graph G , where $\beta = \beta(\epsilon, \kappa)$ is constant whenever κ and ϵ are. The specific dependence of β on κ and ϵ is $\beta(\kappa, \epsilon) = \kappa^{\log \log \kappa - \log \epsilon}$.

3 APPLICATIONS

The result of [4] was used in [3, 5, 6] for computing almost shortest paths in centralized, distributed, streaming, and dynamic centralized models of computations. The basic approach used in these results is to construct a sparse spanner, and then to compute exact shortest paths on the constructed spanner. The sparsity of the latter guarantees that the computation of shortest paths in the spanner is far more efficient than in the original graph.

4 OPEN PROBLEMS

The main open question is whether it is possible to achieve similar results with $\epsilon = 0$. More formally, the question is:

Is it true that for any $\kappa \geq 1$ and any n -vertex graph G there exists $(1, \beta(\kappa))$ -spanner of G with $O(n^{1+1/\kappa})$ edges?

This question was answered in affirmative for κ equal to 2 and 3 [1, 2, 4]. Some lower bounds were recently proved by Woodruff [8].

A less challenging question is to improve the dependence of β on ϵ and κ . Some progress in this direction was achieved by Thorup and Zwick [7].

5 CROSS REFERENCES

None is reported. Entry editors please feel free to add some.

6 RECOMMENDED READING

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- [4] Michael Elkin and David Peleg. $(1 + \epsilon, \beta)$ -spanner constructions for general graphs. In *SIAM J. Comput.*, Vol. 33, No. 3, pp. 608-631, 2004.
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- [6] Liam Roditty and Uri Zwick. Dynamic approximate all-pairs shortest paths in undirected graphs In *Proc. of Symp. on Foundations of Computer Science (FOCS)*, pp.499-508, 2004
- [7] Mikkel Thorup and Uri Zwick. Spanners and Emulators with sublinear distance errors. In *Proc. of Symp. on Discrete Algorithms (SODA)*, pp. 802-809, 2006.
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