

Computer Architecture and Assembly Language

Practical Session 8

Position Independent Code - PIC

- **PIC has everything it needs internally**
- PIC can be placed somewhere in memory, **is executed properly regardless of its absolute address**
- **PIC can be added to any other program,** without fear that it might not work

Position Independent Code - requirements

No direct usage of labels



- No library functions ? ☹️ → only system calls
- One section only
- ‘jump’ and ‘call’ functions ? 😊 → they are relative

No library functions

→ only system calls

We don't know if and where the library functions are.
Thus **there are no “printf” “gets” and so on.**

To perform I/O operation we have to **use the Linux system calls** because INT 0x80 isn't a regular procedure - it is called via the **interrupt table** which **is static.**

One section only

We put all the code in a **single section** – .text (read-only) Or .data (read-write).

Both .text and .data sections **may contain any valid assembly instruction.**

Usage of a single section gives us a possibility to **calculate a relative offset between a pair of code instruction addresses**, and thus **use offset instead of absolute address.**

Only relative jumps

```
1
2          section .data
3 00000000 78563412          numeric:      DD 0x12345678
4 00000004 616263          string:       DB 'abc'
5 00000007 00000000          answer:      DD      0
6
7          section .text
8          global  _start
9
10         _start:
11
12 00000000 60          pushad
13 00000001 6A02          push dword 2
14 00000003 6A01          push dword 1
15 00000005 E815000000    CALL myFunc
16         returnAddress:
17 0000000A A3[07000000]  mov [answer], eax
18 0000000F 83C408          add esp, 8
19 00000012 61          popad
20 00000013 BB00000000    mov ebx,0
21 00000018 B801000000    mov eax,1
22 0000001D CD80          int 0x80
23
24         myFunc:
25 0000001F 55          push  ebp
26 00000020 89E5          mov  ebp, esp
```

If all the code changes its position, **relative jumps are still valid**, because the address difference is preserved.

→ Address of myFunc label = 0x1F
→ Address of the next instruction after the call (i.e. 'mov [answer], eax') is 0xA
→ $0x1F - 0xA = 0x15$, and we get exactly the binary code written here 'E815000000'

No direct usage of labels

```
1
2
3          section .rodata
4 00000000 48656C6C6F0A00          name:  db "Hello",10,0
5
6          global main
7          extern printf
8
9          section .text
10         main:
11 00000000 68[00000000]          push   name
12 00000005 E8(00000000)          call   printf
13 0000000A 83EC04          sub    esp, 4
14 0000000D B801000000          mov    eax, 1
15 00000012 CD80          int   80h
16
```

Labels are resolved at compile time to absolute address

An absolute addresses of name would be resolved based on the offset of name in .rodata section. If the code is moved, the absolute address that was calculated before the moving **would not be correct any more**

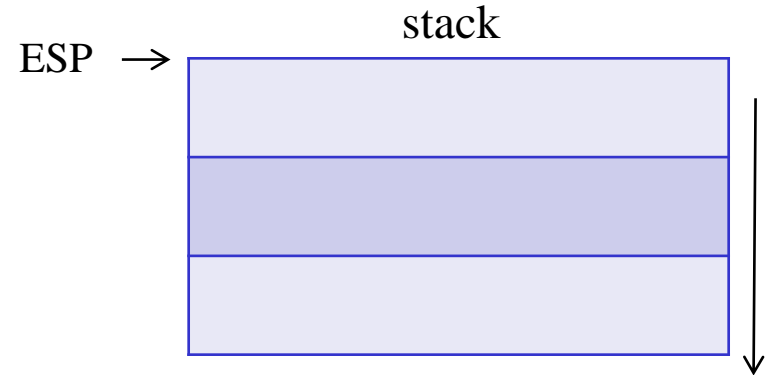
Using labels – PIC example

section .text

```
name:    db      "Hello",10,0
nameLen equ     $ - name

        global _start
get_my_loc:
        call     next_i
next_i:
        pop     edx
        ret

_start:
        call     get_my_loc
        sub     edx, next_i - name
        mov     ecx, edx
        mov     edx, nameLen
        mov     eax, 4
        mov     ebx, 1
        int     80h
        mov     eax, 1
        int     80h
```



Using labels – PIC example

section .text

```
name:    db      "Hello",10,0  
nameLen equ    $ - name
```

```
        global _start  
get_my_loc:
```

```
        call     next_i
```

```
next_i:
```

```
        pop     edx  
        ret
```

```
_start:
```

```
        call    get_my_loc
```

```
        sub     edx, next_i - name
```

```
        mov    ecx, edx
```

```
        mov    edx, nameLen
```

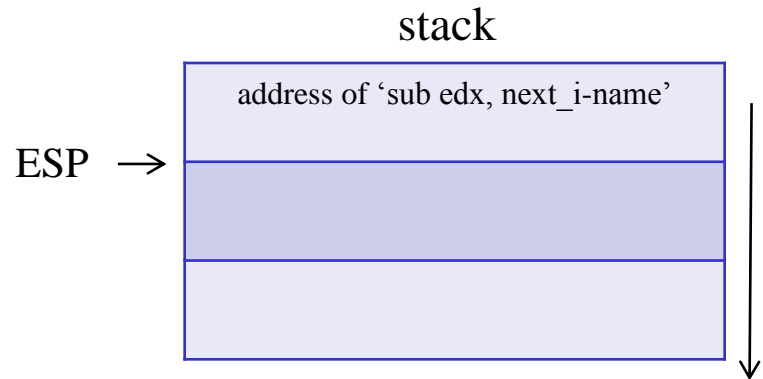
```
        mov    eax, 4
```

```
        mov    ebx, 1
```

```
        int    80h
```

```
        mov    eax, 1
```

```
        int    80h
```



Using labels – PIC example

section .text

```
name:    db      "Hello",10,0  
nameLen equ    $ - name
```

```
global _start
```

```
get_my_loc:
```

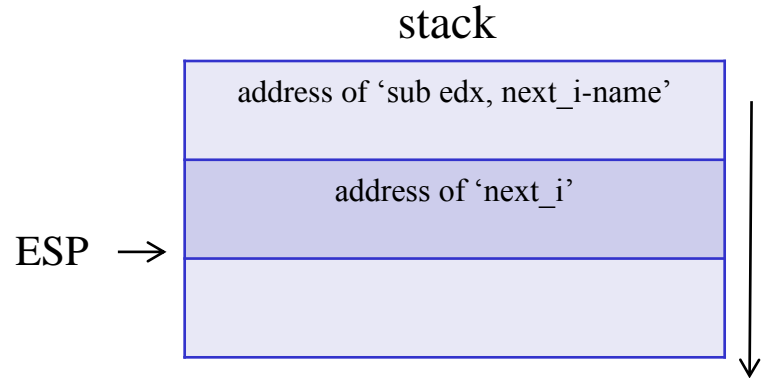
```
call     next_i
```

```
next_i:
```

```
pop     edx  
ret
```

```
_start:
```

```
call     get_my_loc  
sub     edx, next_i - name  
mov     ecx, edx  
mov     edx, nameLen  
mov     eax, 4  
mov     ebx, 1  
int     80h  
mov     eax, 1  
int     80h
```



Using labels – PIC example

section .text

```
name:    db      "Hello",10,0  
nameLen equ    $ - name
```

```
        global _start  
get_my_loc:
```

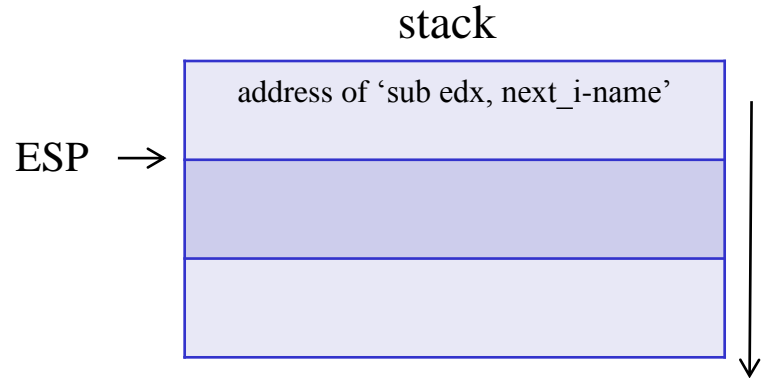
```
        call     next_i
```

```
next_i:
```

```
        pop     edx      ; edx gets address of 'next_i'  
        ret
```

```
_start:
```

```
        call     get_my_loc  
        sub     edx, next_i - name  
        mov     ecx, edx  
        mov     edx, nameLen  
        mov     eax, 4  
        mov     ebx, 1  
        int     80h  
        mov     eax, 1  
        int     80h
```



Using labels – PIC example

section .text

```
name:    db      "Hello",10,0  
nameLen equ    $ - name
```

```
        global _start  
get_my_loc:
```

```
        call     next_i
```

```
next_i:
```

```
        pop     edx
```

```
        ret     ; EIP gets address of 'sub edx, next_i-name'
```

```
_start:
```

```
        call     get_my_loc
```

```
        sub     edx, next_i - name
```

```
        mov     ecx, edx
```

```
        mov     edx, nameLen
```

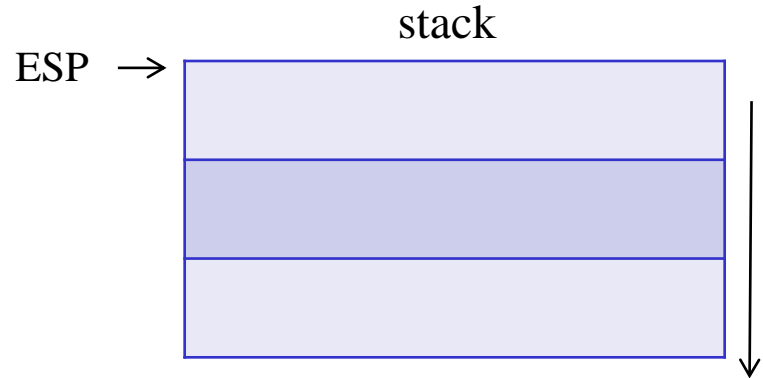
```
        mov     eax, 4
```

```
        mov     ebx, 1
```

```
        int     80h
```

```
        mov     eax, 1
```

```
        int     80h
```



Using labels – PIC example

section .text

```
name:    db        "Hello",10,0  
nameLen equ      $ - name
```

```
    global _start
```

```
get_my_loc:
```

```
    call    next_i
```

```
next_i:
```

```
    pop     edx  
    ret
```

```
_start:
```

```
    call    get_my_loc
```

```
    sub     edx, next_i - name ; edx = 'next_i' address - ('next_i' address - 'name' address)
```

```
    mov     ecx, edx
```

```
    mov     edx, nameLen
```


```
    mov     eax, 4
```

```
    mov     ebx, 1
```

```
    int     80h
```

```
    mov     eax, 1
```

```
    int     80h
```



the address difference between
“next_i” and “name” is constant even
if the code changes it’s position

Using labels – PIC example

section .text

```
name:    db        "Hello",10,0  
nameLen equ      $ - name
```

```
        global _start
```

```
get_my_loc:
```

```
        call       next_i
```

```
next_i:
```

```
        pop        edx
```

```
        ret
```

```
_start:
```

```
        call       get_my_loc
```

```
        sub        edx, next_i - name
```

```
        mov        ecx, edx
```

```
        mov        edx, nameLen ←
```

```
        mov        eax, 4
```

```
        mov        ebx, 1
```

```
        int        80h
```

```
        mov        eax, 1
```

```
        int        80h
```

why we may use 'nameLen' label directly ?

Using labels – PIC example

```
>nasm -f elf sample.s -l sample.lst
```

```
1          section .text
2
3 00000000 48656C6C6F0A00      name:  db      "Hello",10,0
4          nameLen      equ      $ - name
5
6          global _start
7
8          get_my_loc:
9          next_i:
10 0000000C 5A                        pop     edx
11 0000000D C3                        ret
12
13          _start:
14 0000000E E8F4FFFFFF              call   get_my_loc
15 00000013 83EAC                   ← 0x0C = 'next_i' - 'name'      sub     edx, next_i - name
16 00000016 89D1                    mov     ecx, edx
17 00000018 BA07000000              mov     edx, nameLen
18 0000001D B804000000              mov     eax, 4
19 00000022 BB01000000              mov     ebx, 1
20 00000027 CD80                    int     80h
21 00000029 B801000000              mov     eax, 1
22 0000002E CD80                    int     80h
```