

# Compiler Construction

## Exercises

### 1 Review of some Topics in Formal Languages

1.

- (a) Prove that two words  $x, y$  commute (i.e., satisfy  $xy = yx$ ) if and only if there exists a word  $w$  such that  $x = w^m, y = w^n$  for some non-negative integers  $m, n$ .
- (b) Characterize all pairs of words  $x, y, z$  satisfying the equality  $x^2y^2 = z^2$ .
- (c) Let  $k \geq 2$  be an arbitrary fixed integer. Characterize all pairs of words  $x, y, z$  satisfying the equality  $x^ky^k = z^k$ .
- (d) Characterize all triples of words  $x, y, z$  satisfying the equality  $xyz = yzx$ .
- (e) Characterize all triples of words  $x, y, z$  satisfying the equality  $xyz = zyx$ .

2. Let  $\Sigma = \{a, b, \dots, z, 0, 1, \dots, 9, _\cdot\}$ . Let  $L$  be the language consisting of all non-empty words over  $\Sigma$  which (i) do not start with a digit, (ii) do not contain two consecutive occurrences of ‘ $_$ ’, and (iii) do not end with ‘ $_$ ’.

- (a) Construct a regular expression  $r$  such that  $L(r) = L$ .
- (b) Construct a DFA accepting  $L$ .
- (c) How many words of each length  $n$  does  $L$  include?

3. Show that the following languages over  $\Sigma$  are regular:

- (a) The collection of all words whose length is congruent to  $k$  modulo  $m$  (where  $0 \leq k \leq m - 1$ ).
- (b) The collection of all words having the property that, for some fixed positive integer  $k$  and all pairs of letters  $\sigma_1, \sigma_2$ , the difference between the number of occurrences of  $\sigma_1$  and of  $\sigma_2$  in every prefix does not exceed  $k$ .
- (c) The collection of all words containing exactly  $m_1$  occurrences of the word  $w_1$ , exactly  $m_2$  occurrences of the word  $w_2$ , ..., exactly  $m_k$  occurrences of the word  $w_k$ .

4. How many words of length  $n$  do the languages, corresponding to the following regular expressions, contain?

- (a)  $\sigma_1^* \sigma_2^* \dots \sigma_k^*$  (where the  $\sigma_i$ 's are all distinct).
- (b)  $(0 \cup 11 \cup 22 \cup 3333 \cup 4444 \cup 5555 \cup 6666)^*$ .

5. Show that, if  $L$  is a regular language, then so are the following languages:

- (a) The language obtained by replacing, in each word of  $L$ , each occurrence of  $aa$  by  $b$ . (The replacement is done consecutively; thus, the block  $a^{2k}$  is replaced by  $b^k$  and the block  $a^{2k+1}$  by  $b^k a$ .)
- (b) The language obtained from  $L$  by deleting the second last letter in every word of length 2 or more:
- (c) The language consisting of all words in  $L$  in which the number of occurrences of the letter  $\sigma$  is  $r$  modulo  $d$ .
- (d) The language obtained from  $L$  by omitting from each word of  $L$  any number of occurrences of the letter  $\sigma$ . (For example, if  $\sigma_1 \sigma^4 \sigma_2 \sigma^7 \sigma_3 \in L$ , then both words  $\sigma_1 \sigma^2 \sigma_2 \sigma^7 \sigma_3$  and  $\sigma_1 \sigma^3 \sigma_2 \sigma \sigma_3$  belong to the language we construct.)

6. Let  $L_1, L_2$  be two languages over  $\Sigma$ . Show that the “equation”

$$L_1 L \cup L_2 = L$$

has a solution. Moreover, if  $L_1, L_2$  are both regular (or both context-free), then there exists a solution  $L$  with the same property.

7. Let  $\Sigma = \{a, b, c\}$ . Construct DFA's accepting the following languages:

- (a) All words containing neither  $aaa$  nor  $aca$  as a subword.
- (b) All words containing either  $ababa$  or  $abcba$  as a subword.

(c) All words containing both  $a^2$  and  $b^2$ , but not  $c^2$ , as subwords.

**8.** Construct NFAs accepting the languages corresponding to the following regular expressions:

(a)  $bab(bba \cup abb)^*bab$ .

(b)  $ab(ab \cup bba)^* \cup a(ba \cup \phi^*)bba$ .

**9.** Present an algorithm which, given a DFA, returns all words of minimal length accepted by it (or an error if it accepts the empty language).

**10.** Present an algorithm that, given a DFA, returns a DFA accepting a language strictly containing the language accepted by the original DFA and strictly contained in  $\Sigma^*$  (or an error if no such language exists).

**11.**

(a) Show that an infinite regular language may be written as an infinite disjoint union of infinite regular languages.

(b) Show that an infinite context-free language may be written as an infinite disjoint union of infinite context-free languages.

(c) Does an infinite context-free language necessarily contain an infinite regular language?

(d) Does an infinite language, accepted by a Turing machine, necessarily contain an infinite context-free language?

**12.** Show that the following languages are not regular:

(a)  $\{a^m b^n c^{m+n} : m, n \geq 0\}$ .

(b)  $\{a^k b^l c^m d^n : k, l, m, n \geq 0, |\{k, l, m, n\}| \geq 2\}$ .

(c)  $\{0^{m^3+n^3} : m, n \geq 0\}$ .

(c)  $\{0^{l^2+m^3+n^7} : l, m, n \geq 0\}$ .

**13.** Given a set of non-negative integers, the set of their expansions in base 10 forms a partial language of  $\{0, 1, \dots, 9\}^*$ . For each of the following sets show that the corresponding language is regular or not (as indicated):

(a) All powers of 1000 (regular).

(b) All powers of 7 (not regular).

- (c) All perfect cubes (not regular).
- (d)  $\{n! : n \geq 0\}$  (not regular).
- (e) All numbers whose distance from some number of the form  $777 \dots 7$  is at most 7 (regular).

14. Find the languages accepted by the grammars:

(a)

$$S \rightarrow AB \mid BA,$$

$$A \rightarrow aAb \mid \varepsilon,$$

$$B \rightarrow bBa \mid \varepsilon.$$

(b)

$$S \rightarrow \varepsilon,$$

$$S \rightarrow \alpha_i S \alpha_j, \quad 1 \leq i, j \leq m$$

(where  $\Sigma = \{\alpha_1, \alpha_2, \dots, \alpha_m\}$ ).

15. Let  $a_1, a_2, \dots, a_r$  be positive integers and  $b_1, b_2, \dots, b_r$  non-negative integers. Consider the language  $\{\sigma_1^{a_1 n + b_1} \sigma_2^{a_2 n + b_2} \dots \sigma_r^{a_r n + b_r} : n \geq 0\}$ , where  $\sigma_1, \sigma_2, \dots, \sigma_r$  are any letters, not necessarily distinct. Specify the conditions under which this language is context-free.

16. Construct DFA's accepting the same languages as the grammars:

(a)

$$S \rightarrow abS \mid baS \mid bcA,$$

$$A \rightarrow babB,$$

$$B \rightarrow bcbaC,$$

$$C \rightarrow abaC \mid bC \mid \varepsilon.$$

(b)

$$S \rightarrow acbS \mid bcbS \mid bA \mid aC \mid aca,$$

$$A \rightarrow bS \mid caS \mid baB \mid babC \mid a^2,$$

$$B \rightarrow bbA \mid a^2B \mid a^2 \mid b^2 \mid c,$$

$$C \rightarrow bcS \mid aA \mid bB \mid cC \mid cba \mid \varepsilon.$$

17. Construct pushdown automata accepting the same language as the grammars:

(a)

$$S \rightarrow \varepsilon \mid SbS \mid AbS \mid SaB,$$

$$A \rightarrow Bb \mid a^2,$$

$$B \rightarrow b^2S \mid baA \mid b^2.$$

(b)

$$S \rightarrow SaBS \mid bAS \mid ba^2,$$

$$A \rightarrow BaAa \mid aSB \mid ab \mid \varepsilon,$$

$$B \rightarrow b^3B \mid ASA \mid baba.$$

18. Consider the grammar  $G$  defined by:

$$S \rightarrow S + S \mid S * S \mid a \mid b \mid c.$$

(a) Is the language  $L(G)$  regular? If yes – find a DFA accepting the same language as  $G$  and a regular grammar  $G'$  equivalent to  $G$ . If not – prove it.

(b) How many words of each length does the language  $L(G)$  include?

(c) How many derivation trees yield each word in  $L(G)$ ? (Hint: Find a recurrence for the sequence which expresses the required number as a function of the word's length.)

(d) How do your answers to the first three parts change if we add the productions  $S \rightarrow S - S$  and  $S \rightarrow S/S$ ?

## 2 Lexical Analysis

19.

(a) Write a regular expression  $r$  such that  $L(r)$  consists of all identifiers (i.e., all strings of letters and digits, starting with a letter), with the exception of the three strings “if”, “int” and “integer”.

(b) Construct a DFA which recognizes the string “if” as a token of type IF, the strings “int” and “integer” as tokens of type NUM, and all other identifiers as tokens of type ID.

**20.**

- (a) Given an arbitrary fixed integer  $d \geq 2$ , construct a regular expression  $r$  such that  $L(r)$  consists of all base  $d$  expansions of even (positive) integers. (Thus, the alphabet consists of all digits in base  $d$ .)
- (b) Construct a DFA which recognizes every non-negative number, expanded in base  $d$ , as EVEN or as ODD.

**21.**

- (a) Given a regular expression  $r$ , define a regular expression  $\vec{r}$  for vectors (of any non-negative length) of elements of type  $r$ . The entries of a vector are separated by commas and grouped by parentheses.
- (b) Given a DFA recognizing elements of type  $r$ , construct a DFA which recognizes vectors of such elements.
- (c) Can you design your DFA so as to recognize vectors whose length is (i) 3 modulo 10? (ii) a prime?
- (d) The regular expression  $\vec{\vec{r}}$  represents vectors of vectors of elements of type  $r$ . Can you construct a regular expression  $[r]$  for matrices (i.e., vectors whose entries are vectors of the same length) of elements of type  $r$ ?

**22.**

- (a) For any  $n \geq 0$ , write a regular definition for the language of balanced parentheses of nesting level up to  $n$ .
- (b) Write a computer program which, for given  $n$ , will output a DFA for this language. The DFA should inform of the nesting level. (Represent the DFA in any way you like – by a transition table, a graph, etc.)

**23.** Write a regular definition for the language of all strings over  $\{a, b, \dots, z\}$ , not containing “if” as a substring.

**24.** In Java, the command

```
a = b+++--c;
```

passes compilation, whereas the command

```
a = b+++++c;
```

does not. Why?

**25.** In a certain computer language, identifiers are strings of letters and digits, starting with a letter, with the additional constraint that

a character should not appear more than once in the name. (A lower-case letter and the corresponding upper-case letter are considered as distinct.) Construct a DFA, with a minimal possible number of states, recognizing identifiers. How many states does this DFA consist of?

### 3 Syntactic Analysis

**26.** Consider the grammar  $G_1$ , given by:

$$E \rightarrow E + P \mid P,$$

$$P \rightarrow P * V \mid V,$$

$$V \rightarrow a \mid b \mid c.$$

(a) Show that  $L(G_1) = L(G)$ , where  $G$  is the grammar defined in Question 18.

(b) Show that  $G_1$  is unambiguous.

**27.** Consider the grammar  $G_1$  given by:

$$S \rightarrow iSeS \mid iS \mid \varepsilon,$$

and the grammar  $G_2$  given by:

$$S \rightarrow M \mid U,$$

$$M \rightarrow iMeM \mid \varepsilon,$$

$$U \rightarrow iMeU \mid iS.$$

(Intuitively, you should think of these grammars as the two grammars presented in class for conditional statements. Here we deal only with occurrences of the words *if* and *else*, represented by  $i$  and  $e$ , respectively.  $M$  and  $U$  stand for *matched* and *unmatched*, respectively.)

(a) Prove that  $L(G_1) = L(G_2)$ .

(b) Which words are obtained by a unique derivation tree in  $G_1$ ?

(c) Write a program that, given a word in  $\{i, e\}^*$ , finds the number of derivation trees (if any) over  $G_1$  producing this word. Prove that your algorithm works in polynomial time in the length of the input.

(d) Prove that  $G_2$  is unambiguous.

- (e) Write a program that, given a word in  $\{i, e\}^*$ , finds the unique derivation sequence over  $G_2$  producing this word (and gives an error message if the word does not belong to  $L(G_2)$ ).

**28.** Consider the following grammar, designed to solve the ambiguity problem of the **if-then-else** grammar presented in class:

$$stmt \rightarrow \mathbf{if} \ cond \ \mathbf{then} \ stmt \ | \ matched,$$

$$matched \rightarrow \mathbf{if} \ cond \ \mathbf{then} \ matched \ \mathbf{else} \ stmt \ | \ unconditionalStmt.$$

Show that the grammar is still ambiguous.

**29.** Consider the grammar  $G$  given by:

$$S \rightarrow SS\sigma_1 \ | \ SS\sigma_2 \ | \ \dots \ | \ SS\sigma_r \ | \ \sigma_{r+1},$$

where  $\sigma_1, \sigma_2, \dots, \sigma_{r+1}$  are distinct terminals. Is the grammar unambiguous? If yes – prove your claim, if not – produce a word in  $L(G)$  with two distinct derivation trees.

**30.** Consider the grammar  $G$ , given by:

$$S \rightarrow FS \ | \ \varepsilon,$$

$$F \rightarrow aB \ | \ bA,$$

$$A \rightarrow a \ | \ bAA,$$

$$B \rightarrow b \ | \ aBB.$$

- (a) Show that  $L(G)$  consists of all words over  $\{a, b\}$  with an equal number of occurrences of  $a$  and of  $b$ .
- (b) Show that  $G$  is unambiguous.

**31.** Construct an unambiguous grammar  $G$ , such that  $L(G)$  consists of all words over  $\{a, b\}$  in which the number of occurrences of  $a$  is not less than that of  $b$ .

**32.** For each of the following grammars, find the minimal  $k$  for which it is  $LL(k)$ . In case no such  $k$  exists, determine whether the grammar is unambiguous.

(a)

$$E \rightarrow M + E \ | \ M - E \ | \ M * E \ | \ M/E \ | \ M,$$

$$M \rightarrow V \ | \ V ++ \ | \ V --,$$

$$V \rightarrow a \ | \ b \ | \ c.$$



(b)

$$\begin{aligned} E &\rightarrow E + M \mid E - M \mid E * M \mid E/M \mid M, \\ M &\rightarrow V \mid V ++ \mid V --, \\ V &\rightarrow a \mid b \mid c. \end{aligned}$$

(c)

$$\begin{aligned} E &\rightarrow M + E \mid M - E \mid M * E \mid M/E \mid M, \\ M &\rightarrow V \mid V ++ \mid V -- \mid ++V \mid --V, \\ V &\rightarrow a \mid b \mid c. \end{aligned}$$

**33.** Two of the algorithms discussed in class enable us deciding which letters  $X \in N \cup T$  have the property that there exists a word  $w \in L$  such that  $X \xRightarrow{*} w$ , where  $L$  is any one of the two languages  $T^*$  and  $\{\varepsilon\}$ . Show that you can do the same for any regular language  $L \subseteq T^*$ . (You may use algorithms studied in the Automata course without detailing them.)

**34.** Find the FIRST sets of all non-terminals and right-hand sides of all rules for the following grammars:

(a)

$$\begin{aligned} S &\rightarrow ABc, \\ A &\rightarrow a \mid \varepsilon, \\ B &\rightarrow b \mid \varepsilon. \end{aligned}$$

(b)

$$\begin{aligned} S &\rightarrow aS \mid AS \mid BAb, \\ A &\rightarrow Aab \mid AB \mid \varepsilon, \\ B &\rightarrow Aa \mid BbB \mid \varepsilon. \end{aligned}$$

(c)

$$\begin{aligned} S &\rightarrow aSe \mid A, \\ A &\rightarrow bAe \mid B, \\ B &\rightarrow cBe \mid d. \end{aligned}$$

(d)

$$\begin{aligned} S &\rightarrow ABCS \mid SS \mid aba, \\ A &\rightarrow ACB \mid cb \mid \varepsilon, \\ B &\rightarrow BCB \mid A \mid bc, \\ C &\rightarrow AS \mid c. \end{aligned}$$

(e)

$$\begin{aligned}S &\rightarrow SS \mid AB \mid c, \\A &\rightarrow Aa \mid Aab \mid a \mid \varepsilon, \\B &\rightarrow bC \mid bB \mid Sb \mid b, \\C &\rightarrow cA \mid SC \mid c.\end{aligned}$$

**35.** Find the FOLLOW sets of all non-terminals for the grammars in Question 34.

**36.** Determine whether each of the following grammars is  $LR(k)$  for some  $k$ . If yes – find the minimal such  $k$ . (Hint: In some cases it may be helpful to find first the minimal  $k$  for which the grammar is  $LL(k)$ , and then use the fact that the grammar is  $LR(k)$  for this  $k$ .)

(a)  $S \rightarrow a^2Sb^3 \mid a^3b^4$ .

(b)  $S \rightarrow aSa \mid \varepsilon$ .

(c)  $S \rightarrow aSba \mid aba$ .

(d) The grammar defined in Question 18.

(e) The grammar defined in Question 26.

(f) The grammar defined in Question 29.

**37.**

(a) Find the left contexts of all non-terminals and the  $LR(0)$  contexts of all rules for the grammars in the preceding question.

(b) Same for the two grammars in Question 27.