

Mayer Goldberg Numbers

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Abstract

In this paper we analyze the so-called *Mayer Goldberg numbers* which contain all possible finite information encoded in given bases.

1 Introduction

Consider the following binary number:

$$\text{mg}(2) = 0.10001101100000101001110010111011100000001 \dots_2$$

What's special about it? Taking a closer look, we discover its underlying structure:

$$\text{mg}(2) = \underbrace{0.1}_{\text{chunk 1}} \underbrace{00\ 01\ 10\ 11}_{\text{chunk 2}} \underbrace{000\ 001\ 010\ 011\ 100\ 101\ 110\ 111}_{\text{chunk 3}} \dots_2 \quad (1.1)$$

This number, when encoded in binary, is a concatenation of all words over $\{0, 1\}^+$ in increasing order of their value in base-2. Thus, it “contains” within itself binary encodings of all possible deterministic structures.

2 Computation

In order to represent $\text{mg}(2)$ algebraically, we shall compute it via summation. For this purpose the following formula will prove helpful:

$$\sum_{k=1}^m kx^k = \frac{(mx - m - 1)x^{m+1} + x}{(x - 1)^2} \quad (2.1)$$

It immediately follows from derivation of $\sum_{k=0}^m x^k$ by x .

Let us count the digits in $\text{mg}(2)$ from left to right, starting with 0. Then, the index of the leftmost digit in chunk $n \geq 1$ is:

$$\text{offs}(n) = \sum_{k=1}^{n-1} k2^k = (n - 2)2^n + 2 \quad (2.2)$$

Now $mg(2)$ can be rewritten as

$$\begin{aligned}
 mg(2) &= 2^{-\text{offs}(1)} \cdot \underline{0.1} \\
 &\quad + 2^{-\text{offs}(2)} \cdot \underline{0.0 \underline{01} \underline{10} \underline{11}} \\
 &\quad + 2^{-\text{offs}(3)} \cdot \underline{0.00 \underline{001} \underline{010} \underline{011} \underline{100} \underline{101} \underline{110} \underline{111}} \\
 &\quad \dots \\
 &= \sum_{n=1}^{\infty} \frac{\text{chunk}(n)}{2^{\text{offs}(n)}} \\
 &= \frac{1}{4} \sum_{n=1}^{\infty} \frac{\text{chunk}(n)}{2^{(n-2)2^n}}
 \end{aligned} \tag{2.3}$$

Where

$$\begin{aligned}
 \text{chunk}(n) &= \frac{1}{2^{n-1}} \sum_{k=1}^{2^n-1} k 2^{-nk} \\
 &= \frac{1}{2^{n-1}} \sum_{k=1}^{2^n-1} k (2^{-n})^k \\
 &= \frac{(2^{-n})^{2^n} ((2^n - 1)2^{-n} - 2^n) + 2^{-n}}{(2^{-n} - 1)^2 2^{n-1}} \quad \text{by (2.1)} \\
 &= 2 \cdot \frac{2^{-n} - (2^n + 2^{-n} - 1)2^{-n2^n}}{2^n + 2^{-n} - 2}
 \end{aligned} \tag{2.4}$$

Finally resulting in

$$\boxed{mg(2) = \frac{1}{2} \sum_{n=1}^{\infty} \frac{2^{-n} - (2^n + 2^{-n} - 1)2^{-n2^n}}{(2^n + 2^{-n} - 2)2^{(n-2)2^n}}} \tag{2.5}$$

First 51 digits of $mg(2)$ in decimal representation:

$$\boxed{mg(2) = 0.55277423455897304746839735242380246111017997632137\dots}$$

3 Generalization

Mayer Goldberg number for base-2 can be easily generalized for other bases:

$$\boxed{mg(\beta) = \beta^{1 - \frac{\beta}{(\beta-1)^2}} \sum_{n=1}^{\infty} \frac{\beta^{-n} - (\beta^n + \beta^{-n} - 1)\beta^{-n\beta^n}}{(\beta^n + \beta^{-n} - 2)\beta^{\frac{((n-1)\beta-n)\beta^n}{(\beta-1)^2}}}} \tag{3.1}$$

For example, the number for base-4, corresponding to all possible information encoded on a 4-hole punch-tape is:

$$\boxed{mg(4) = 0.421944444444444444444444353739145534438965151883420879\dots}$$

Note that $mg(3) > mg(2)$.

4 Further Exploration

Here we shall ask how many such numbers are — how many numbers, when represented in base β , possess the property that they “contain” all words in $\{0, \dots, \beta - 1\}^+$.

Easy to see, for each base $\beta \geq 2$ there are \aleph_1 such numbers. This can be shown by transposing words in the chunks of $\text{mg}(\beta)$, for example.

However, in my opinion it is very unlikely that most of the numbers maintain this property.¹ I don't even think that this property is shared over bases in general, except for special cases like $\beta_1 = k\beta_2$.

¹Like they do with probability of 1, for example, in respect to equal representation of all digits in any base.