

The MST of Symmetric Disk Graphs (in Arbitrary Metrics) is Light

Shay Solomon

Department of Computer Science,
Ben-Gurion University of the Negev
P.O. Box 653, Beer-Sheva 84105, ISRAEL
E-mail: shayso@cs.bgu.ac.il
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Abstract

Consider an n -point metric $M = (V, \delta)$, and a transmission range assignment $r : V \rightarrow \mathbb{R}^+$ that maps each point $v \in V$ to the disk of radius $r(v)$ around it. The *symmetric disk graph* (henceforth, SDG) that corresponds to M and r is the undirected graph over V whose edge set includes an edge (u, v) if both $r(u)$ and $r(v)$ are no smaller than $\delta(u, v)$. SDGs are often used to model wireless communication networks.

Abu-Affash, Aschner, Carmi and Katz (SWAT 2010, [1]) showed that for any *2-dimensional Euclidean* n -point metric M , the weight of the MST of every *connected* SDG for M is $O(\log n) \cdot w(\text{MST}(M))$, and that this bound is tight. However, the upper bound proof of [1] relies heavily on basic geometric properties of 2-dimensional Euclidean metrics, and does not extend to higher dimensions. A natural question that arises is whether this surprising upper bound of [1] can be generalized for wider families of metrics, such as 3-dimensional Euclidean metrics.

In this paper we generalize the upper bound of Abu-Affash et al. [1] for Euclidean metrics of any dimension. Furthermore, our upper bound extends to *arbitrary metrics* and, in particular, it applies to any of the normed spaces ℓ_p . Specifically, we demonstrate that for *any* n -point metric M , the weight of the MST of every connected SDG for M is $O(\log n) \cdot w(\text{MST}(M))$.

*Department of Computer Science, Ben-Gurion University of the Negev, POB 653, Beer-Sheva 84105, Israel.

E-mail: {shayso}@cs.bgu.ac.il

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1 Introduction

1.1 The MST of Symmetric Disk Graphs. Consider a network that is represented as an (undirected) weighted graph $G = (V, E, w)$, and assume that we want to compute a spanning tree for G of *small weight*, i.e., of weight that is close to the weight $w(MST(G))$ of the minimum spanning tree (MST) for G . (See Section 1.6 for the definition of weight.) However, due to some physical constraints (e.g., network faults) we are only given a connected spanning subgraph G' of G , rather than G itself. In this situation it is natural to use the MST of the given subgraph G' . The *weight-coefficient* of G' with respect to G is defined as the ratio between $w(MST(G'))$ and $w(MST(G))$. If the weight-coefficient of G' is small enough, we can use $MST(G')$ as a spanning tree for G of small weight.

The problem of computing spanning trees of small weight (especially the MST) is a fundamental one in Computer Science [20, 18, 7, 27, 14, 10], and the above scenario arises naturally in many practical contexts (see, e.g., [31, 13, 36, 24, 25, 26, 11, 12]). In particular, this scenario is motivated by wireless network design.

In this paper we focus on the *symmetric disk graph model* in wireless communication networks, which has been subject to considerable research. (See [17, 15, 16, 23, 32, 5, 34, 1], and the references therein.) Let $M = (V, \delta)$ be an n -point metric that is represented as a complete weighted graph $G(M) = (V, \binom{V}{2}, w)$ in which the weight $w(e)$ of each edge $e = (u, v)$ is equal to $\delta(u, v)$. Also, let $r : V \rightarrow \mathbb{R}^+$ be a transmission range assignment that maps each point $v \in V$ to the disk of radius $r(v)$ around it. The *symmetric disk graph* (henceforth, SDG) that corresponds to M and r , denoted $SDG(M, r)$,¹ is the undirected spanning subgraph of $G(M)$ whose edge set includes an edge $e = (u, v)$ if both $r(u)$ and $r(v)$ are no smaller than $w(e)$. Under the symmetric disk graph model we cannot use all the edges of $G(M)$, but rather only those that are present in $SDG(M, r)$. Clearly, if $r(v) \geq diam(M)^2$ for each point $v \in V$, then $SDG(M, r)$ is simply the complete graph $G(M)$. However, the transmission ranges are usually significantly shorter than $diam(M)$, and many edges that belong to $G(M)$ may not be present in $SDG(M, r)$. Therefore, it is generally impossible to use the MST of M under the symmetric disk graph model, simply because some of the edges of $MST(M)$ are not present in $SDG(M, r)$ and thus cannot be accessed. Instead, assuming the weight-coefficient of $SDG(M, r)$ with respect to M is small enough, we can use $MST(SDG(M, r))$ as a spanning tree for M of small weight.

Abu-Affash et al. [1] showed that for any *2-dimensional Euclidean* n -point metric M , the weight of the MST of every *connected* SDG for M is $O(\log n) \cdot w(MST(M))$. In other words, they proved that for any 2-dimensional Euclidean n -point metric, the weight-coefficient of every *connected* SDG is $O(\log n)$. In addition, Abu-Affash et al. [1] provided a matching lower bound of $\Omega(\log n)$ on the weight-coefficient of connected SDGs that applies to a basic 1-dimensional Euclidean metric. Notably, the upper bound proof of [1] relies heavily on basic geometric properties of 2-dimensional Euclidean metrics, and does not extend to higher dimensions. A natural question that arises is whether the logarithmic upper bound of [1] on the weight-coefficient of connected SDGs can be generalized for wider families of metrics, such as 3-dimensional Euclidean metrics.

¹The definition of symmetric disk graph can be generalized in the obvious way for any weighted graph. Specifically, the *symmetric disk graph* $SDG(G, r)$ that corresponds to a weighted graph $G = (V, E, w)$ and a range assignment r is the undirected spanning subgraph of G whose edge set includes an edge $e = (u, v) \in E$ if both $r(u)$ and $r(v)$ are no smaller than $w(e)$.

²The *diameter* of a metric M , denoted $diam(M)$, is defined as the largest pairwise distance in M .

In this paper we generalize the upper bound of Abu-Affash et al. [1] for Euclidean metrics of any dimension. Furthermore, our upper bound extends to *arbitrary metrics* and, in particular, it applies to any of the normed spaces ℓ_p . Specifically, we demonstrate that for *any* n -point metric M , every connected SDG has weight-coefficient $O(\log n)$. In fact, our upper bound is even more general, applying to unconnected SDGs as well. That is, we show that the weight of the minimum spanning forest (MSF) of every (possibly unconnected) SDG for M is $O(\log n) \cdot w(MST(M))$.

The fact that the weight-coefficient of SDGs for arbitrary metrics is relatively small is quite surprising. In particular, we demonstrate that for other basic parameters of spanning trees, the situation is fundamentally different. Consider, for example, the maximum degree (henceforth, degree) parameter. Clearly, for any metric there is a spanning tree with degree 2. On the other hand, consider an n -point metric M^* in which the distance between a designated point $rt \in M^*$ and every other point is equal to 1, and all other distances are equal to 2. The SDG corresponding to M^* and the range assignment $r \equiv 1$ that maps each point to the unit disk around it is the n -star graph rooted at rt , having degree $n - 1$. Thus, the *degree-coefficient* of connected SDGs for metrics can be as large as $\Omega(n)$ in general, which is *exponentially larger* than the weight-coefficient. (See Section 4.2 for the formal definition of degree-coefficient.) We show that the same lower bound of $\Omega(n)$ also applies to other parameters of spanning trees, including radius and depth, diameter and hop-diameter, sum of all pairwise distances, and sum of all distances from a designated vertex.

Finally, we remark that our logarithmic upper bound on the weight-coefficient of SDGs does not extend to general (undirected) weighted graphs. Indeed, if the weight function of the graph does not satisfy the triangle inequality, the weight-coefficient of SDGs can be arbitrarily large even for complete graphs. Moreover, we demonstrate that there are (non-complete) 1-dimensional Euclidean n -vertex graphs³ for which the weight-coefficient of SDGs can be as large as $\Omega(n)$.

1.2 The Range Assignment Problem. Given a network $G = (V, E, w)$, a *range assignment* for G is an assignment of transmission ranges to each of the vertices of G . A range assignment is called *complete* if the induced (directed) communication graph is strongly connected. In the *range assignment problem* the objective is to find a complete range assignment for which the total power consumption (henceforth, cost) is minimized. The power consumed by a vertex $v \in V$ is $r(v)^\alpha$, where $r(v)$ is the range assigned to v and $\alpha \geq 1$ is some constant. Thus the cost of the range assignment is given by $\sum_{v \in V} r(v)^\alpha$. The range assignment problem was first studied by Kirousis et al. [19], who proved that the problem is NP-hard in 3-dimensional Euclidean metrics, assuming $\alpha = 2$, and also presented a simple 2-approximation algorithm. Subsequently, Clementi et al. [9] proved that the problem remains NP-hard in 2-dimensional Euclidean metrics.

We believe that it is more realistic to study the range assignment problem under the symmetric disk graph model. Specifically, the potential transmission range of a vertex v is bounded by some maximum range $r(v)$, and any two vertices u, v can directly communicate with each other if and only if v lies within the range assigned to u and vice versa. Blough et al. [3] showed that this version of the range assignment problem is also NP-hard in 2-dimensional and 3-dimensional Euclidean metrics. Also, Calinescu et al. [4] devised a $(1 + \frac{1}{2} \ln 3 + \epsilon)$ -approximation scheme and a more practical $(\frac{15}{8})$ -approximation algorithm. Abu-Affash et al. [1] showed that, assuming $\alpha = 1$, the cost of an optimal range assignment with bounds on the ranges is greater by at most a logarithmic factor than the cost of an optimal range assignment without such bounds. This result of Abu-

³A *1-dimensional Euclidean graph* is a weighted graph in which the vertices represent points on a line, and the weight of each edge is equal to the Euclidean distance between its endpoints.

Affash et al. [1] is a simple corollary of their upper bound on the weight-coefficient of SDGs for 2-dimensional Euclidean metrics. Consequently, this result of [1] for the range assignment problem holds only in 2-dimensional Euclidean metrics. By applying our generalized upper bound on the weight-coefficient of SDGs, we extend this result of Abu-Affash et al. [1] to arbitrary metrics.

1.3 Proof Overview. As was mentioned above, the upper bound proof of [1] is very specific, and relies heavily on basic geometric properties of 2-dimensional Euclidean metrics. Hence, it does not apply to 3-dimensional Euclidean metrics, let alone to arbitrary metrics. Our upper bound proof is based on completely different principles. In particular, it is independent of the geometry of the metric and applies to every complete graph whose weight function satisfies the triangle inequality. In fact, at the heart of our proof is a lemma that applies to an even wider family of graphs, namely, the family of all traceable⁴ weighted graphs. Specifically, let S and H be an SDG and a minimum-weight Hamiltonian path of some traceable weighted n -vertex graph G , respectively, and let F be the MSF of S . Our lemma states that there is a set \tilde{E} of edges in F of weight at most $w(H)$, such that the graph $F \setminus \tilde{E}$ obtained by removing all edges of \tilde{E} from F contains at least $\frac{1}{5} \cdot n$ isolated vertices. The proof of this lemma is based on a delicate combinatorial argument that does not assume either that the graph G is complete or that its weight function satisfies the triangle inequality. We believe that this lemma is of independent interest. (See Lemma 2.2 in Section 2.) By employing this lemma inductively, we are able to show that the weight of F is bounded above by $\log_{\frac{5}{4}} n \cdot w(H)$, which, by the triangle inequality, yields an upper bound of $2 \cdot \log_{\frac{5}{4}} n$ on the weight-coefficient of S with respect to G . Interestingly, our upper bound of $2 \cdot \log_{\frac{5}{4}} n$ on the weight-coefficient of SDGs for arbitrary metrics improves the corresponding upper bound of [1] (namely, $90 \cdot \log_{\frac{5}{4}} n + 1$), which holds only in 2-dimensional Euclidean metrics, by a multiplicative factor of 45.

1.4 Related Work on Disk Graphs. The symmetric disk graph model is a generalization of the extremely well-studied *unit disk graph model* (see, e.g., [8, 24, 21, 26, 22]). The *unit disk graph* of a metric M , denoted $UDG(M)$, is the symmetric disk graph corresponding to M and the range assignment $r \equiv 1$ that maps each point to the unit disk around it. (It is usually assumed that M is a 2-dimensional Euclidean metric.) It is easy to see that in the case when $UDG(M)$ is connected, all edges of $MST(M)$ belong to $UDG(M)$, and so $MST(UDG(M)) = MST(M)$. Hence the weight-coefficient of connected unit disk graphs for arbitrary metrics is equal to 1. In the general case, we note that all edges of $MSF(UDG(M))$ belong to $MST(M)$, and so the weight-coefficient of (possibly unconnected) unit disk graphs for arbitrary metrics is at most 1.

Another model that has received much attention in the literature is the *asymmetric disk graph model* (see, e.g., [21, 33, 29, 30, 1]). The *asymmetric disk graph* corresponding to a metric $M = (V, \delta)$ and a range assignment r is the directed graph over V , where there is an arc of weight $\delta(u, v)$ from u to v if $r(u) \geq \delta(u, v)$. On the negative side, Abu-Affash et al. [1] provided a lower bound of $\Omega(n)$ on the weight-coefficient of strongly connected asymmetric disk graphs that applies to a 1-dimensional Euclidean n -point metric. However, asymmetric communication models are generally considered to be impractical, because in such models many communication primitives become unacceptably complicated [28, 35]. In particular, the asymmetric disk graph model is often viewed as less realistic than the symmetric disk graph model, where, as was mentioned above, we obtain a logarithmic upper bound on the weight-coefficient for arbitrary metrics.

⁴A graph is called *traceable* if it contains a Hamiltonian path.

1.5 Structure of the Paper. The main result of this paper is given in Section 2. Therein we obtain a logarithmic upper bound on the weight-coefficient of SDGs for arbitrary metrics. In Section 3 we provide an application of this upper bound to the range assignment problem. Finally, Section 4 is devoted to negative results concerning two possible extensions of the upper bound of Section 2. We start (Section 4.1) with showing that this upper bound does not extend to general graphs, and proceed (Section 4.2) showing that in comparison to the weight parameter, other basic parameters of spanning trees incur significantly larger bounds.

1.6 Preliminaries. Given a (possibly weighted) graph G , its vertex set (respectively, edge set) is denoted by $V(G)$ (resp., $E(G)$). For an edge set $E' \subseteq E(G)$, we denote by $G \setminus E'$ the graph obtained by removing all edges of E' from G . Similarly, for an edge set E'' over the vertex set $V(G)$, we denote by $G \cup E''$ the graph obtained by adding all edges of E'' to G . The weight of an edge e in G is denoted by $w(e)$. For an edge set $E \subseteq E(G)$, its weight $w(E)$ is defined as the sum of all edge weights in E , i.e., $w(E) = \sum_{e \in E} w(e)$. The weight of G is defined as the weight of its edge set $E(G)$, namely, $w(G) = w(E(G))$. Finally, for a positive integer n , we denote the set $\{1, 2, \dots, n\}$ by $[n]$.

2 The MST of SDGs is Light

In this section we prove that the weight-coefficient of SDGs for arbitrary n -point metrics is $O(\log n)$.

We will use the following well-known fact in the sequel.

Fact 2.1 *Let G be a weighted graph in which all edge weights are distinct. Then G has a unique MSF, and the edge of maximum weight in every cycle of G does not belong to the MSF of G .*

In what follows we assume for simplicity that all the distances in any metric are distinct. This assumption does not lose generality, since any ties can be broken using, e.g., lexicographic rules. We may henceforth assume that there is a unique MST for any metric, and a unique MSF for every SDG of any metric.

The following lemma is central in our upper bound proof.

Lemma 2.2 *Let $M = (V, \delta)$ be an n -point metric and let $r : V \rightarrow \mathbb{R}^+$ be a range assignment. Also, let $F = (V, E_F)$ be the MSF of the symmetric disk graph $S = \text{SDG}(M, r)$ and let $H = (V, E_H)$ be a minimum-weight Hamiltonian path of M . Then there is an edge set $\tilde{E} \subseteq E_F$ of weight at most $w(H)$, such that the graph $F \setminus \tilde{E}$ contains at least $\frac{1}{5} \cdot n$ isolated vertices.*

Remark: This statement remains valid if instead of the metric M we take a general traceable weighted graph.

Proof: Denote by E' the set of edges in H that belong to the SDG S , i.e., $E' = E_H \cap E(S)$, and let $E'' = E_H \setminus E'$ be the complementary edge set of E' in E_H . Also, denote by E'_1 the set of edges in E' that belong to the MSF F , i.e., $E'_1 = E' \cap E_F$, and let $E'_2 = E' \setminus E'_1$ be the complementary edge set of E'_1 in E' . Note that (1) $E' \subseteq E(S)$, (2) $E'' \cap E(S) = \emptyset$, (3) $E'_1 \subseteq E_F$, and (4) $E'_2 \cap E_F = \emptyset$.

Write $F_0 = F$, $k = |E'_2|$, and let e'_1, e'_2, \dots, e'_k denote the edges of E'_2 by increasing order of weight. Next, we construct k spanning forests F_1, F_2, \dots, F_k of S in the following iterative process. For each index $i = 1, 2, \dots, k$, the graph $F_{i-1} \cup \{e'_i\}$ obtained from F_{i-1} by adding to it the edge

e'_i contains a unique cycle C_i . Since H is cycle-free, at least one edge of C_i does not belong to H . Let \tilde{e}_i be an arbitrary such edge, and denote by $F_i = F_{i-1} \cup \{e'_i\} \setminus \{\tilde{e}_i\}$ the graph obtained from F_{i-1} by adding to it the edge e'_i and removing the edge \tilde{e}_i .

It is easy to verify that the cycles C_1, C_2, \dots, C_k that are identified during this process are subgraphs of the symmetric disk graph S . Moreover, for each index $i \in [k]$, we have $E(C_i) \subseteq E_F \cup E'_2$, yielding $\tilde{e}_i \in E_F \setminus E_H$.

Write $\tilde{E}_2 = \{\tilde{e}_1, \tilde{e}_2, \dots, \tilde{e}_k\}$, and notice that $|\tilde{E}_2| = |E'_2|$, $\tilde{E}_2 \subseteq E_F \setminus E_H$. The following claim implies that $w(\tilde{E}_2) \leq w(E'_2)$.

Claim 2.3 *For each index $i \in [k]$, $w(\tilde{e}_i) \leq w(e'_i)$.*

Proof: Fix an arbitrary index $i \in [k]$, and define $E'_{(i)} = \{e'_1, \dots, e'_i\}$. Recall that the cycle C_i is a subgraph of S , and notice that each edge of C_i that do not belong to F must belong to $E'_{(i)}$, i.e., $E(C_i) \setminus E_F \subseteq E'_{(i)}$. Fact 2.1 implies that the edge of maximum weight in C_i , denoted e_i^* , does not belong to F , and so $e_i^* \in E'_{(i)}$. Since e'_i is the edge of maximum weight in $E'_{(i)}$, it holds that $w(e_i^*) \leq w(e'_i)$. Also, as \tilde{e}_i belongs to C_i , we have by definition $w(\tilde{e}_i) \leq w(e_i^*)$. Claim 2.3 follows. \square

Denote by $H'' = H \setminus E' = H \setminus (E'_1 \cup E'_2)$ and $F'' = F \setminus (E'_1 \cup \tilde{E}_2)$ the graphs obtained from H and F by removing all edges of $E' = E'_1 \cup E'_2$ and $E'_1 \cup \tilde{E}_2$, respectively. By definition, $E(H'') = E''$. For an edge $e = (u, v)$, denote by $\min(e)$ the endpoint of e with smaller radius, i.e., $\min(e) = u$ if $r(u) < r(v)$, and $\min(e) = v$ otherwise. Consider an arbitrary edge $e \in E''$. Since no edge of E'' belongs to the symmetric disk graph S , it follows that $r(\min(e)) < w(e)$. Also, since the graph F'' is a subgraph of S , the weight of every edge that is incident to $\min(e)$ in F'' is no greater than $r(\min(e)) < w(e)$.

Next, we remove some edges of the graphs H'' and F'' and add them to the two initially empty edge sets E_H^* and E_F^* , respectively. This is done in the following way. We initialize $E_H^* = E_F^* = \emptyset$, and then examine the edges of E'' one after another in an arbitrary order. For each edge $e \in E''$, we check whether the vertex $\min(e)$ is isolated in F'' or not. If $\min(e)$ is isolated in F'' , we leave H'', F'', E_H^* , and E_F^* intact. Otherwise, at least one edge is incident to $\min(e)$ in F'' . Let \tilde{e} be an arbitrary such edge, and note that $w(\tilde{e}) \leq r(\min(e)) < w(e)$. We remove the edge e from the graph H'' and add it to the edge set E_H^* , and remove the edge \tilde{e} from the graph F'' and add it to the edge set E_F^* . This process is repeated iteratively until all edges of E'' have been examined. Notice that at each stage of this process, it holds that $|E_F^*| = |E_H^*|$, $E_F^* \subseteq E_F \setminus E_H$, $w(E_F^*) \leq w(E_H^*)$.

In what follows we consider the graphs H'', F'' and edge sets E_H^*, E_F^* resulting at the end of this process. Define $\tilde{E} = E'_1 \cup \tilde{E}_2 \cup E_F^*$. By construction, we have $\tilde{E} \subseteq E_F$, $F'' = F \setminus \tilde{E}$. The following claim completes the proof of Lemma 2.2.

Claim 2.4 (1) $w(\tilde{E}) \leq w(H)$. (2) *The graph $F'' = F \setminus \tilde{E}$ contains at least $\frac{1}{5} \cdot n$ isolated vertices.*

Proof: We start with proving the first assertion of the claim. Define $E = E' \cup E_H^* = E'_1 \cup E'_2 \cup E_H^*$. It is easy to see that the edge sets E'_1, E'_2 , and E_H^* are pairwise disjoint subsets of E_H , yielding $w(E'_1) + w(E'_2) + w(E_H^*) = w(E) \leq w(H)$. Recall that $w(\tilde{E}_2) \leq w(E'_2)$ and $w(E_F^*) \leq w(E_H^*)$. Consequently,

$$\begin{aligned} w(\tilde{E}) &= w\left(E'_1 \cup \tilde{E}_2 \cup E_F^*\right) \leq w(E'_1) + w(\tilde{E}_2) + w(E_F^*) \\ &\leq w(E'_1) + w(E'_2) + w(E_H^*) = w(E) \leq w(H). \end{aligned}$$

Next, we prove the second assertion of the claim. Denote by m_H (respectively, m_F) the number $|E(H'')|$ (resp., $|E(F'')|$) of edges in the graph H'' (resp., F'').

Suppose first that $m_F < \frac{2}{5} \cdot n$. Observe that in any n -vertex graph with m edges there are at least $n - 2m$ isolated vertices. Thus, the number of isolated vertices in F'' is bounded below by $n - 2m_F > n - \frac{4}{5} \cdot n = \frac{1}{5} \cdot n$, as required.

We henceforth assume that $m_F \geq \frac{2}{5} \cdot n$.

Recall that $|\tilde{E}_2| = |E'_2|, |E_F^*| = |E_H^*|$. It is easy to see that the edge sets $E'_1, \tilde{E}_2, E_F^*, E'_2$, and E_H^* are pairwise disjoint, and so

$$|\tilde{E}| = |E'_1| + |\tilde{E}_2| + |E_F^*| = |E'_1| + |E'_2| + |E_H^*| = |E|.$$

Observe that $H'' = H \setminus E$ and recall that $F'' = F \setminus \tilde{E}$, yielding $|E(H'')| = |E_H| - |E|, |E(F'')| = |E_F| - |\tilde{E}|$. Also, note that $|E_H| = n - 1 \geq |E_F|$. Altogether,

$$m_H = |E(H'')| = |E_H| - |E| \geq |E_F| - |\tilde{E}| = |E(F'')| = m_F. \quad (1)$$

Next, observe that for each edge e in H'' , the vertex $\min(e)$ is isolated in F'' . By definition, for any pair e, e' of non-incident edges in H'' , $\min(e) \neq \min(e')$. Since the graph H'' is cycle-free and the maximum degree of a vertex in H'' is at most two, it follows that the number m_H of edges in H'' is at most twice greater than the number of isolated vertices in F'' . Thus, the number of isolated vertices in F'' is bounded below by $\frac{1}{2} \cdot m_H \geq \frac{1}{2} \cdot m_F \geq \frac{1}{5} \cdot n$. (The first inequality follows from (1) whereas the second inequality follows from the above assumption.)

This completes the proof of claim 2.4. □

Lemma 2.2 follows. □

Next, we employ Lemma 2.2 inductively to upper bound the weight of SDGs in terms of the weight of the minimum-weight Hamiltonian path of the metric. The desired upper bound of $O(\log n)$ on the weight-coefficient of SDGs for arbitrary n -point metrics would immediately follow.

Lemma 2.5 *Let $M = (V, \delta)$ be an n -point metric and let $r : V \rightarrow \mathbb{R}^+$ be a range assignment. Also, let $F = (V, E_F)$ be the MSF of the symmetric disk graph $S = \text{SDG}(M, r)$ and let $H = (V, E_H)$ be a minimum-weight Hamiltonian path of M . Then $w(F) \leq \log_{\frac{5}{4}} n \cdot w(H)$.*

Proof: The proof is by induction on the number n of points in the metric M .

Basis: $n \leq 4$. The case $n = 1$ is trivial. Suppose next that $2 \leq n \leq 4$. In this case $\log_{\frac{5}{4}} n \geq \log_{\frac{5}{4}} 2 > 3$. Also, the MSF F of S contains at most 3 edges. By the triangle inequality, the weight of each edge of F is bounded above by the weight $w(H)$ of the Hamiltonian path H . Hence, $w(F) \leq 3 \cdot w(H) < \log_{\frac{5}{4}} n \cdot w(H)$.

Induction step: We assume that the statement holds for all smaller values of n , $n \geq 5$, and prove it for n . By Lemma 2.2, there is an edge set $\tilde{E} \subseteq E_F$ of weight at most $w(H)$, such that the set I of isolated vertices in the graph $F \setminus \tilde{E}$ satisfies $|I| \geq \frac{1}{5} \cdot n$. Consider the complementary edge set $\hat{E} = E_F \setminus \tilde{E}$ of edges in F . Observe that no edge of \hat{E} is incident to a vertex of I . Denote by \hat{M} the sub-metric of M induced by the point set of $\hat{V} = V \setminus I$, let $\hat{S} = \text{SDG}(\hat{M}, r)$ be the SDG corresponding to \hat{M} and the original range assignment r , and let $\hat{F} = (\hat{V}, E_{\hat{F}})$ be the MSF of \hat{S} . Notice that the induced subgraph of S over the vertex set \hat{V} is equal to \hat{S} , and so \hat{E} is a subgraph of \hat{S} . Since \hat{F} is a spanning forest of \hat{S} , replacing the edge set \tilde{E} of F by the edge

set $E_{\hat{F}}$ does not affect the connectivity of the graph, i.e., the graph $\bar{F} = F \setminus \hat{E} \cup E_{\hat{F}}$ that is obtained from F by removing the edge set \hat{E} and adding the edge set $E_{\hat{F}}$ has exactly the same connected components as F . Thus, by breaking all cycles in the graph \bar{F} , we get a spanning forest of S . The weight of this spanning forest is bounded above by the weight $w(\bar{F}) = w(F \setminus \hat{E} \cup E_{\hat{F}})$ of the graph \bar{F} , and is bounded below by the weight $w(F)$ of the MSF F of S . It follows that $w(\hat{E}) \leq w(E_{\hat{F}}) = w(\hat{F})$. Write $\hat{n} = |\hat{V}|$, and let $\hat{H} = (\hat{V}, E_{\hat{H}})$ be a minimum-weight Hamiltonian path of \hat{M} . Since $|I| \geq \frac{1}{5} \cdot n$, we have

$$\hat{n} = |\hat{V}| = |V \setminus I| \leq \frac{4}{5} \cdot n \leq n - 1.$$

(The last inequality holds for $n \geq 5$.) By the induction hypothesis for \hat{n} , $w(\hat{F}) \leq \log_{\frac{5}{4}} \hat{n} \cdot w(\hat{H})$. Also, the triangle inequality implies that $w(\hat{H}) \leq w(H)$. Hence,

$$\begin{aligned} w(\hat{E}) &\leq w(\hat{F}) \leq \log_{\frac{5}{4}} \hat{n} \cdot w(\hat{H}) \leq \log_{\frac{5}{4}} \left(\frac{4}{5} \cdot n \right) \cdot w(H) \\ &= \log_{\frac{5}{4}} n \cdot w(H) - w(H). \end{aligned}$$

We conclude that

$$\begin{aligned} w(F) &= w(E_F) = w(\tilde{E}) + w(E_F \setminus \tilde{E}) = w(\tilde{E}) + w(\hat{E}) \\ &\leq w(H) + \log_{\frac{5}{4}} n \cdot w(H) - w(H) = \log_{\frac{5}{4}} n \cdot w(H). \end{aligned}$$

□

By the triangle inequality, the weight of the minimum-weight Hamiltonian path of any metric is at most twice greater than the weight of the MST for that metric. We derive the main result of this paper.

Theorem 2.6 *For any n -point metric $M = (V, \delta)$ and any range assignment $r : V \rightarrow \mathbb{R}^+$, $w(\text{MSF}(\text{SDG}(M, r))) = O(\log n) \cdot w(\text{MST}(M))$.*

3 The Range Assignment Problem

In this section we demonstrate that for any metric, the cost of an optimal range assignment with bounds on the ranges is greater by at most a logarithmic factor than the cost of an optimal range assignment without such bounds. This result follows as a simple corollary of the upper bound given in Theorem 2.6.

Let $M = (V, \delta)$ be an n -point metric, and assume that the n points of V , denoted by v_1, v_2, \dots, v_n , represent transceivers. Also, let $r' : V \rightarrow \mathbb{R}^+$ be a function that provides a maximum transmission range for each of the points of V . In the *bounded range assignment problem* the objective is to compute a range assignment $r : V \rightarrow \mathbb{R}^+$, such that (i) for each point $v_i \in V$, $r(v_i) \leq r'(v_i)$, (ii) the induced SDG (using the ranges $r(v_1), r(v_2), \dots, r(v_n)$), namely $\text{SDG}(M, r)$, is connected, and (iii) $\sum_{i=1}^n r(v_i)$ is minimized. In the *unbounded range assignment problem* the maximum transmission range for each of the points of V is unbounded, i.e., $r'(v_i) = \text{Diam}(M)$,

for each point $v_i \in V$. The function r' is called a *bounding function*. Also, the sum $\sum_{i=1}^n r(v_i)$ is called the *cost* of the range assignment r , and is denoted by $COST(r)$.

Fix an arbitrary bounding function $r' : V \rightarrow \mathbb{R}^+$. Denote by $OPT(M, r')$ the cost of an optimal solution for the bounded range assignment problem corresponding to M and r' . Also, denote by $OPT(M)$ the cost of an optimal solution for the unbounded range assignment problem corresponding to M . Clearly, $OPT(M, r') \geq OPT(M)$. Next, we show that $OPT(M, r') = O(\log n) \cdot OPT(M)$.

Let $SDG(M, r')$ be the SDG corresponding to M and r' , and let T be the MST of $SDG(M, r')$. We define r to be the range assignment that assigns $r(v_i)$ with the weight of the heaviest edge incident to v_i in T , for each point $v_i \in V$. By construction, $r(v_i) \leq r'(v_i)$, for each point $v_i \in V$. Also, notice that the SDG corresponding to M and r , namely $SDG(M, r)$, contains T and is thus connected. Hence, the range assignment r provides a feasible solution for the bounded range assignment problem corresponding to M and r' , yielding $OPT(M, r') \leq COST(r)$. By a double counting argument, it follows that $\sum_{i=1}^n r(v_i) \leq 2 \cdot w(T)$. Also, by Theorem 2.6, $w(T) = w(MST(SDG(M, r'))) = O(\log n) \cdot w(MST(M))$. Finally, it is easy to verify that $w(MST(M)) \leq OPT(M)$. Altogether,

$$\begin{aligned} OPT(M, r') &\leq COST(r) = \sum_{i=1}^n r(v_i) \leq 2 \cdot w(T) \\ &= O(\log n) \cdot w(MST(M)) = O(\log n) \cdot OPT(M). \end{aligned}$$

4 Negative Results

4.1 The MST of SDGs in General Graphs is Heavy

In this section we show that the upper bound of Theorem 2.6 on the weight-coefficient of SDGs does not extend to general (undirected) weighted graphs.

Let W be an arbitrary large number.

Suppose first that the weight function of the graph does not satisfy the triangle inequality. Let C_3 be the 3-cycle graph on the vertex set $\{a, b, c\}$ in which the edge (a, b) has unit weight, the edge (a, c) has weight 2, and the edge (b, c) has weight W . Observe that the SDG corresponding to C_3 and the range assignment r that assigns $r(a) = 1, r(b) = r(c) = W$, denoted $SDG(C_3, r)$, contains the two edges (a, b) and (b, c) but does not contain the edge (a, c) . Thus, $SDG(C_3, r)$ is a spanning tree of C_3 of weight $W + 1$. On the other hand, the MST of the original graph C_3 consists of the two edges (a, b) and (a, c) and has weight 3. It follows that the weight-coefficient of $SDG(C_3, r)$ with respect to C_3 is $\Omega(W)$. In other words, the weight-coefficient of SDGs can be arbitrary large in general.

When the weight function of the graph satisfies the triangle inequality, the weight of each edge in the MST (or the MSF) of the SDG is bounded above by the weight of the entire MST of the original graph. Hence, the weight-coefficient of SDGs in this case is at most linear in the number of vertices in the graph.

Next, we provide a matching lower bound that applies to (non-complete) 1-dimensional Euclidean graphs. Let $\epsilon > 0$ be a tiny number, and consider the set V of n points $s, u_1, u_2, u_3, \dots, u_{n-2}, t$ that lie on the x -axis with coordinates $0, 1, 1 + \epsilon, 1 + 2\epsilon, \dots, 1 + (n - 3) \cdot \epsilon, W + 1$, respectively,

where $W \gg n$ and $\epsilon \ll \frac{1}{n}$. Define $U = V \setminus \{s, t\} = \{u_1, u_2, \dots, u_{n-2}\}$, and observe that the Euclidean distance between s (respectively, t) and each point of $U \setminus \{u_1\}$ is slightly larger than 1 (resp., slightly smaller than W), whereas the Euclidean distance between s (respectively, t) and u_1 is equal to 1 (resp., W). Let $G = (V, E, \|\cdot\|)$ be the Euclidean graph over the point set V , where $E = \{(s, u_i) \mid u_i \in U\} \cup \{(u_i, t) \mid u_i \in U\}$. (See Figure 1 for an illustration.) Observe that the

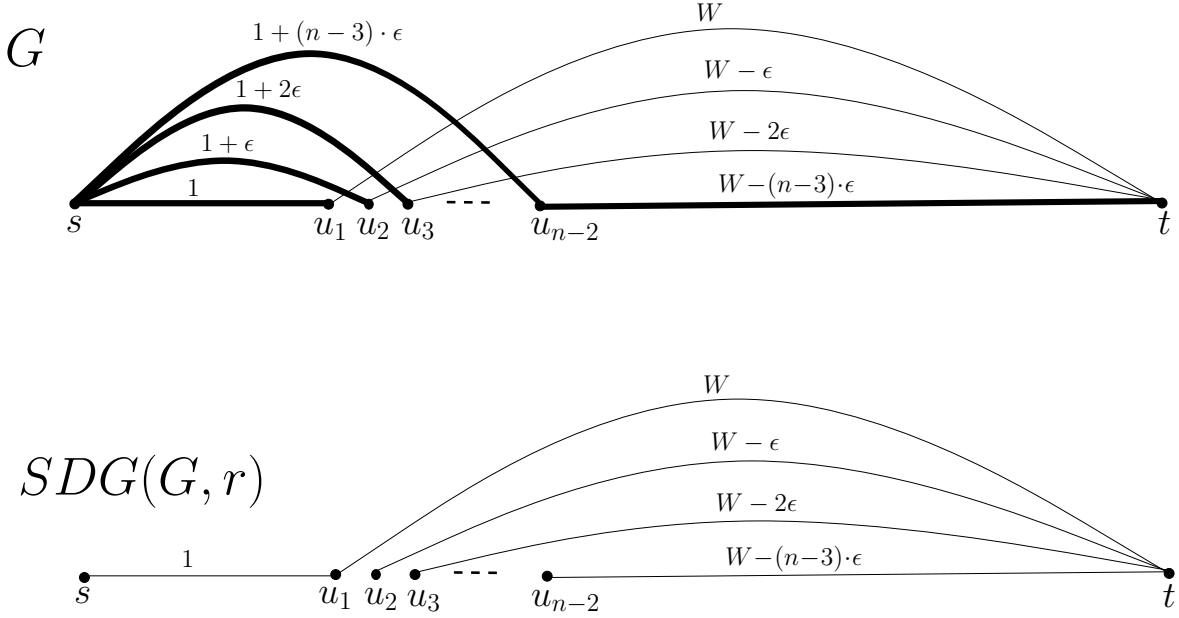


Figure 1: The graph G that appears at the top of the figure is a 1-dimensional Euclidean n -vertex graph that contains $2n - 4$ edges. The $n - 1$ edges of the MST of G are depicted by bold lines. The graph $SDG(G, r)$ that appears at the bottom of the figure is the SDG corresponding to G and r , where $r(s) = 1$, $r(u_1) = r(u_2) = \dots = r(u_{n-2}) = r(t) = W$.

MST of G contains all $n - 2$ edges $\{(s, u_i) \mid u_i \in U\}$ that are incident to s in G and another edge (u_{n-2}, t) that connects t with its closest neighbor u_{n-2} in G . Thus, $w(MST(G)) \approx (n - 2) + W$. Consider the range assignment r that assigns $r(x) = W$, for any point $x \in V \setminus \{s\}$, and $r(s) = 1$. The SDG corresponding to G and r , denoted $SDG(G, r)$, contains all $n - 2$ edges of G that are incident to t and another edge (s, u_1) that connects s with its closest neighbor u_1 in G . In particular, $SDG(G, r)$ is a spanning tree for G of weight roughly $(n - 2) \cdot W + 1$. Thus, the weight-coefficient of $SDG(G, r)$ with respect to G is roughly $\frac{(n-2) \cdot W + 1}{(n-2) + W} \approx n - 2$, assuming $W \gg n$ and $\epsilon \ll \frac{1}{n}$. In other words, we obtained a lower bound of $\Omega(n)$ on the weight-coefficient of SDGs for 1-dimensional Euclidean n -vertex graphs.

4.2 The Weight is Exponentially Better than Other Parameters

Our main result (Theorem 2.6) implies that the weight-coefficient of SDGs for arbitrary n -point metrics is $O(\log n)$. In contrast, we showed (see Section 1.1) that the degree-coefficient of SDGs for n -point metrics can be as large as $\Omega(n)$. (The formal definition of degree-coefficient is given below.) In this section we demonstrate that, similarly to the degree parameter, a lower bound of $\Omega(n)$ also applies to other basic parameters of spanning trees.

Let $T = (T, rt)$ be a (possibly weighted) rooted tree. Before we proceed, we provide the formal definitions of some basic parameters of trees:

- The *maximum degree* (or shortly, *degree*) of T is the maximum number of edges that are incident to some vertex in T .
- The *radius* (respectively, *depth*) of T is the maximum weighted (resp., unweighted) distance between rt and some leaf in T .
- The *diameter* (respectively, *hop-diameter*) of T is the maximum weighted (resp., unweighted) distance between some pair of vertices in T .
- The *sum of all pairwise distances* (or shortly, *sum-pairwise*) of T is defined as the sum of weighted distances between all pairs of vertices in T .
- The *sum of all single-source distances* (or shortly, *sum-single*) of T is defined as the sum of weighted distances between rt and all other vertices in T .

(The sum-pairwise and sum-single parameters also have unweighted versions.)

Let G be a (possibly weighted) graph and let G' be a connected spanning subgraph of G . The *degree-coefficient* of G' with respect to G is defined as the ratio between the degree of the minimum-degree spanning tree of G' and the degree of the minimum-degree spanning tree of G . In exactly the same way we can define the *radius-coefficient*, *depth-coefficient*, *diameter-coefficient*, *hop-diameter-coefficient*, *sum-pairwise-coefficient*, and *sum-single-coefficient*.

In Section 1.1 we showed that the degree-coefficient of SDGs can be as large as $\Omega(n)$ for n -point metrics. Next, we provide a simple example for which each of the other parameters defined above incurs the same lower bound of $\Omega(n)$.

Consider the n -point metric $M = (V, \delta)$, where $V = \{v_1, v_2, \dots, v_n\}$, and the distance function δ satisfies that $\delta(v_1, v_2) = \delta(v_2, v_3) = \dots = \delta(v_{n-1}, v_n) = 1$, and for all other pairs of points $v_i, v_j \in V$, $\delta(v_i, v_j) = 2$.

Let $SDG(M, r)$ be the SDG corresponding to M and the range assignment $r \equiv 1$. It is easy to see that $SDG(M, r)$ is the unweighted n -path (v_1, v_2, \dots, v_n) .

Let $T = (T, v_1)$ be the spanning tree of M rooted at v_1 that consists of the $n - 1$ edges $(v_1, v_2), (v_1, v_3), \dots, (v_1, v_n)$. Thus, T is the n -star graph rooted at v_1 in which the weight of the edge (v_1, v_2) is equal to 1, and all other edge weights are equal to 2. Notice that $SDG(M, r)$ is already a spanning tree of M . Denote by $S = (S, v_1)$ the tree $SDG(M, r)$ rooted at v_1 . It is easy to see that:

- Both the radius and depth of T are $O(1)$, whereas the corresponding measures of S are $\Omega(n)$. Thus, both the radius-coefficient and depth-coefficient of $SDG(M, r)$ with respect to M are $\Omega(n)$.
- Both the diameter and hop-diameter of T are $O(1)$, whereas the corresponding measures of S are $\Omega(n)$. Thus, both the diameter-coefficient and hop-diameter-coefficient of $SDG(M, r)$ with respect to M are $\Omega(n)$.
- The sum-pairwise of T is $O(n^2)$ whereas the sum-pairwise of S is $\Omega(n^3)$. Thus, the sum-pairwise-coefficient of $SDG(M, r)$ with respect to M is $\Omega(n)$.

- The sum-single of T is $O(n)$ whereas the sum-single of S is $\Omega(n^2)$. Thus, the sum-single-coefficient of $SDG(M, r)$ with respect to M is $\Omega(n)$.

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